SPDK, NVME-OF Acceleration

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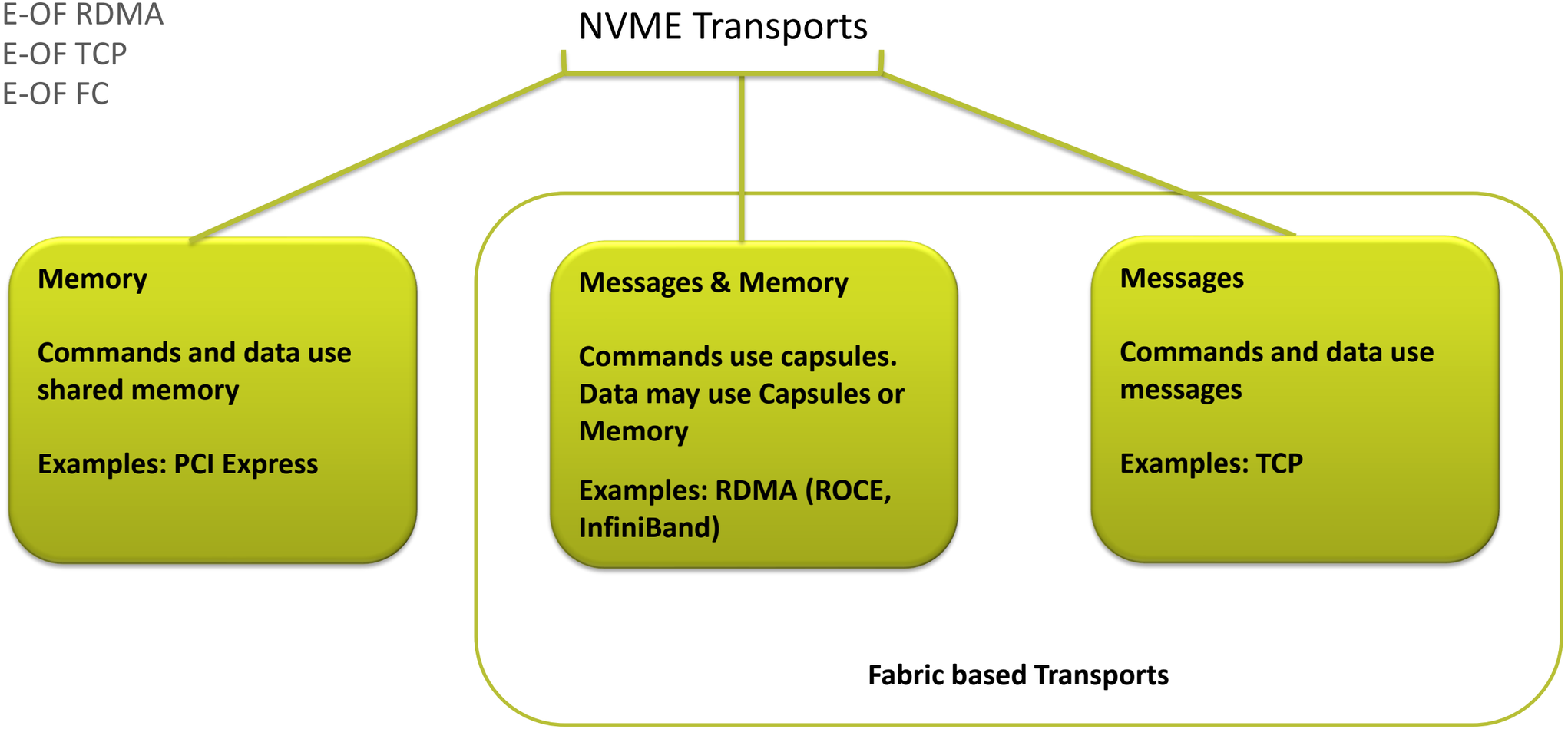
September, 2019

Agenda

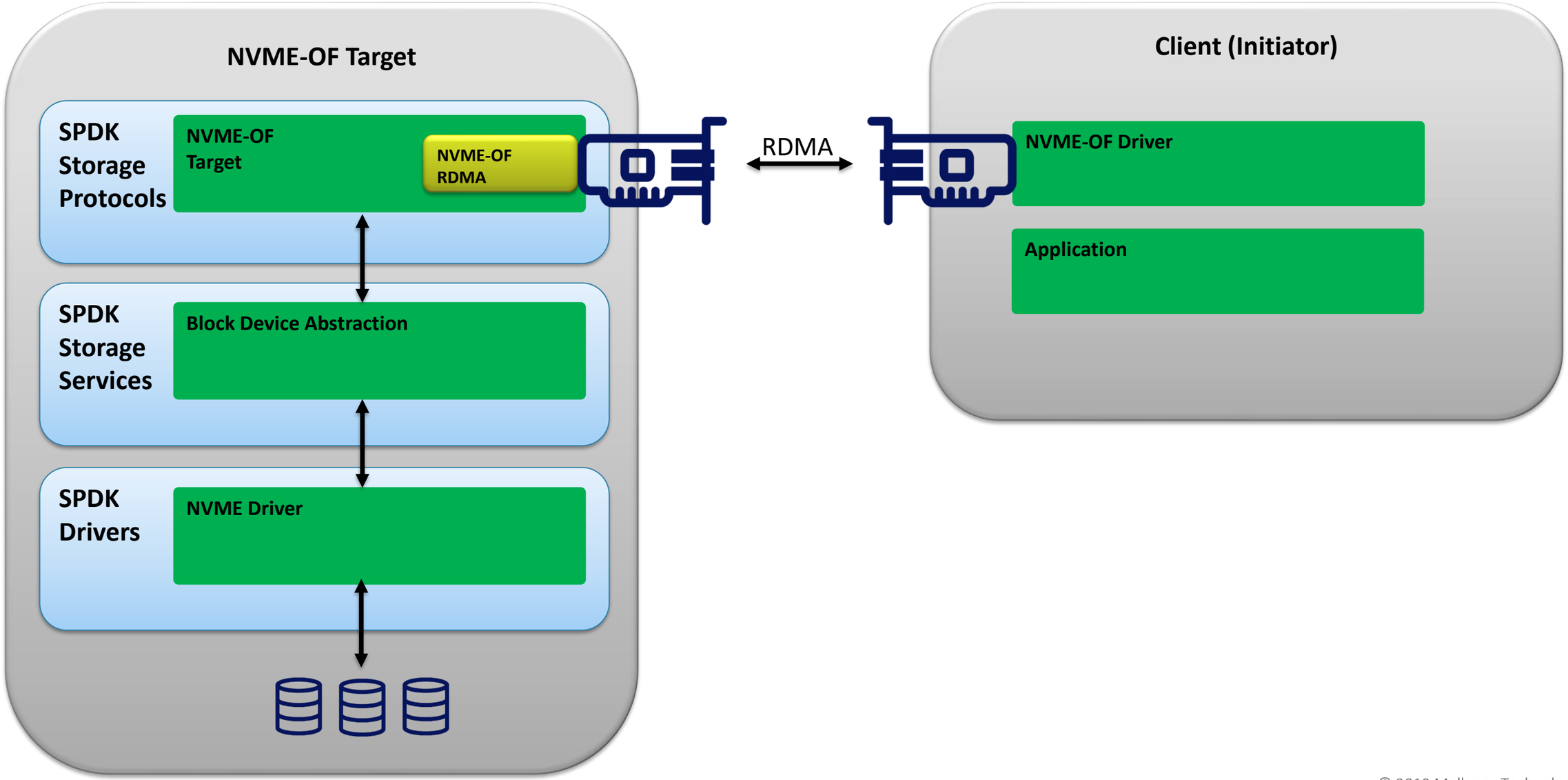
- Background
- Low-level optimizations in NVME-OF RDMA transport
- Data protection in RDMA transport
- Advanced hardware accelerations in network layer

NVME and NVME-OF

- **NVMe is designed to work over a PCIe bus**
- The **NVMe over Fabrics** is the protocol used for transferring NVMe storage commands between the client nodes over storage fabric
  - NVME-OF RDMA
  - NVME-OF TCP
  - NVME-OF FC



SPDK. NVME-OF Abstraction



NVME-OF RDMA Optimizations

NVME-OF RDMA. Performance optimizations

- Scope
  - NVME-OF Target on x86
  - NVME-OF Target on ARM
  - NVME-OF Target forwards IO to backend target
- Network cards
  - “ConnectX-5”
  - “BlueField”

RDMA. Selective signaling

- “Selective signaling” reduces PCIe bandwidth and CPU usage by eliminating DMA completion
- In IO Read flow, RDMA_WRITE is followed by RDMA_SEND
  - Completion for RDMA_WRITE can be skipped
- Developed by Alexey Marchuk, Mellanox: https://review.gerrithub.io/c/spdk/spdk/+/456469
  - Available in SPDK v19.07
- “Selective signaling” increases IOPs in “randread”
  - ARM up to 15%

RDMA. Work request batching

- “Work request batching” reduces CPU use and PCIe bandwidth by using single MMIO operation (“Doorbell”) for multiple requests
- The default approach for WQE (work request element) transferring requires separate MMIO for each WQE
- WQE batching improve:
  - IO Read flow: RDMA_WRITE is followed by RDMA_SEND
  - “Heavy” loads (high queue depth): NVME-OF Target needs to submit multiple RDMA operations
  - Multi –element SGL: Each element needs own RDMA operation
- Developed by:
  - Seth Howell, “Intel”: https://review.gerrithub.io/c/spdk/spdk/+/449265 - NVME-OF Target
    - Available in SPDK v19.07 . Requires applying fix : https://review.gerrithub.io/c/spdk/spdk/+/466029
  - Evgenene Kotchetov, “Mellanox” : https://review.gerrithub.io/c/spdk/spdk/+/462585/ - NVME-OF Initiator
- Preliminary results:
  - ARM: randread (queue depth 64) up to 5%, randwrite (queue depth 64) up to 12% increase in IOPs

RDMA. Work request's payload inlining

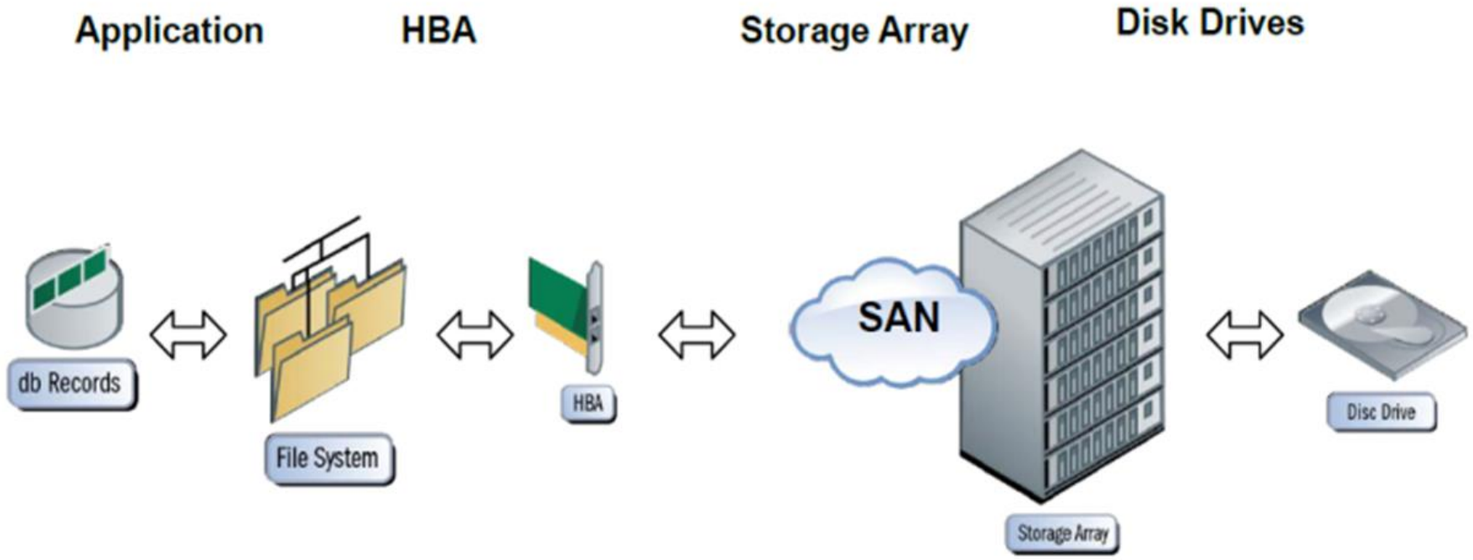
- Payload inlining reduces PCIe bandwidth by eliminating DMA read for payload
- Small payloads up to a few hundred of bytes can be encapsulated into WQE
- Payload inlining can be used for NVME-OF response
  - Capsule size is 16 bytes
  - The feature is under development (Alexey Marchuk, "Mellanox"): https://github.com/Mellanox/spdk/commit/8682d067e5ab9470fb3596db0c47411c974ac47f

NVME-OF RDMA
Data protection

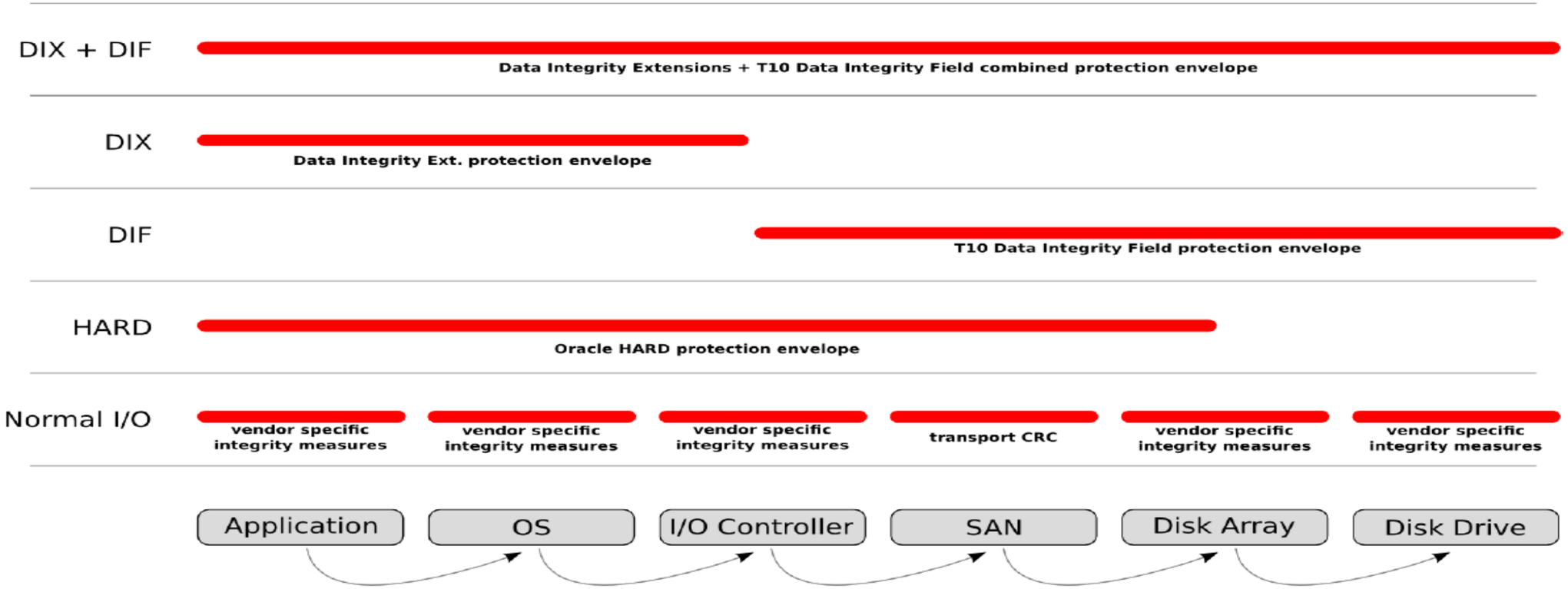
Data Corruption

- IS A DISSASTER !!
- Backups may have bad data
- Downtime/Corruptions may be fatal to a company
- It is better to Not Return any data, than return a wrong one
- Occur as a result of bugs, both SW and HW (drivers, HBAs, Disks, Arrays)
- Common failures:
  - Write incorrect data to the storage device – may take months for recognition
  - Misdirected writes

Error can happen in every entity in the IO path:

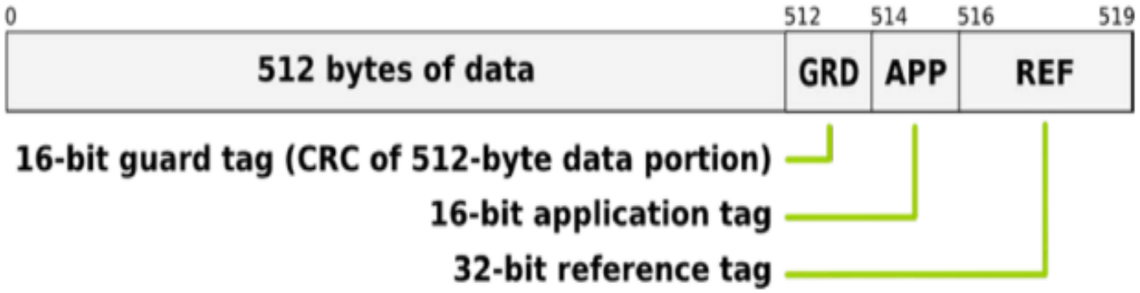


I/O path entities



*based on Martin K. Peterson slide

Model

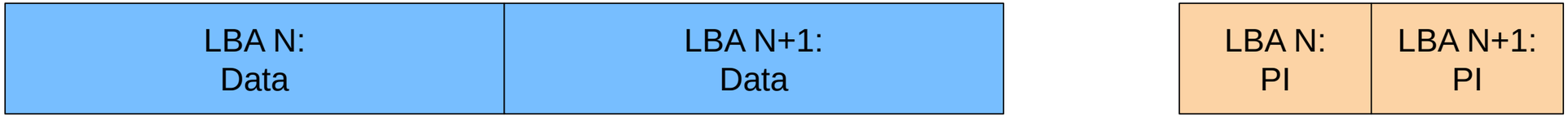


- 8 byte of integrity tuple per sector
- Guard tag:
  - Per request property
  - Protects the data portion of the sector
  - On the Wire – CRC using well-defined polynomial
  - OS – usually use cheaper IP checksum algorithm (may use CRC)
  - I/O controller should convert between types, if needed
- Application tag:
  - Opaque
  - Free usage by application
- Reference tag:
  - Protect against misdirected writes
  - Type 1 - 32 LSbits of the LBA are used as base tag and incremented with each segment
  - Type 2 - 32 LSbits of the LBA used as base tag, can be anything for the rest
  - Type 3 – Only Guard tag is checked

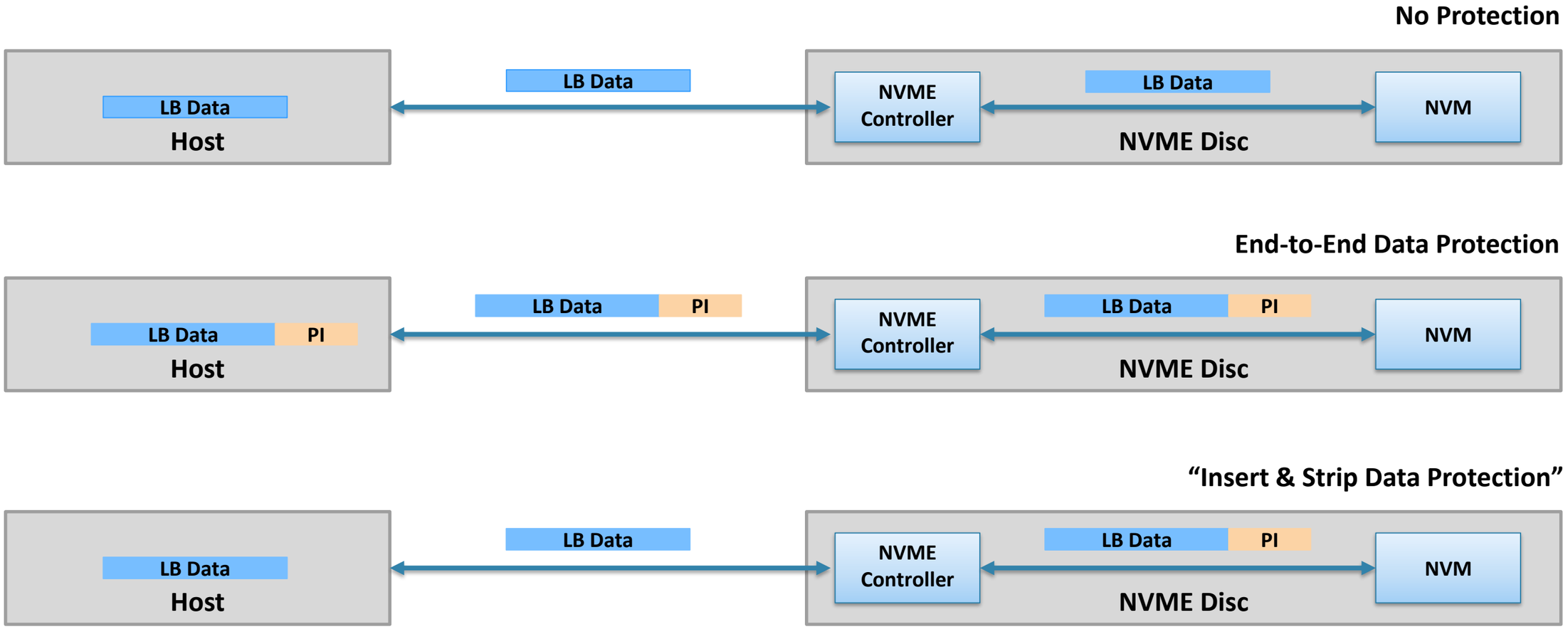
NVMEoF – Metadata Handling

- Two possibilities for MetaData layout
  - Interleaved: Each data block is appended with 8byte integrity payload.
  - Not supported by Linux for local devices

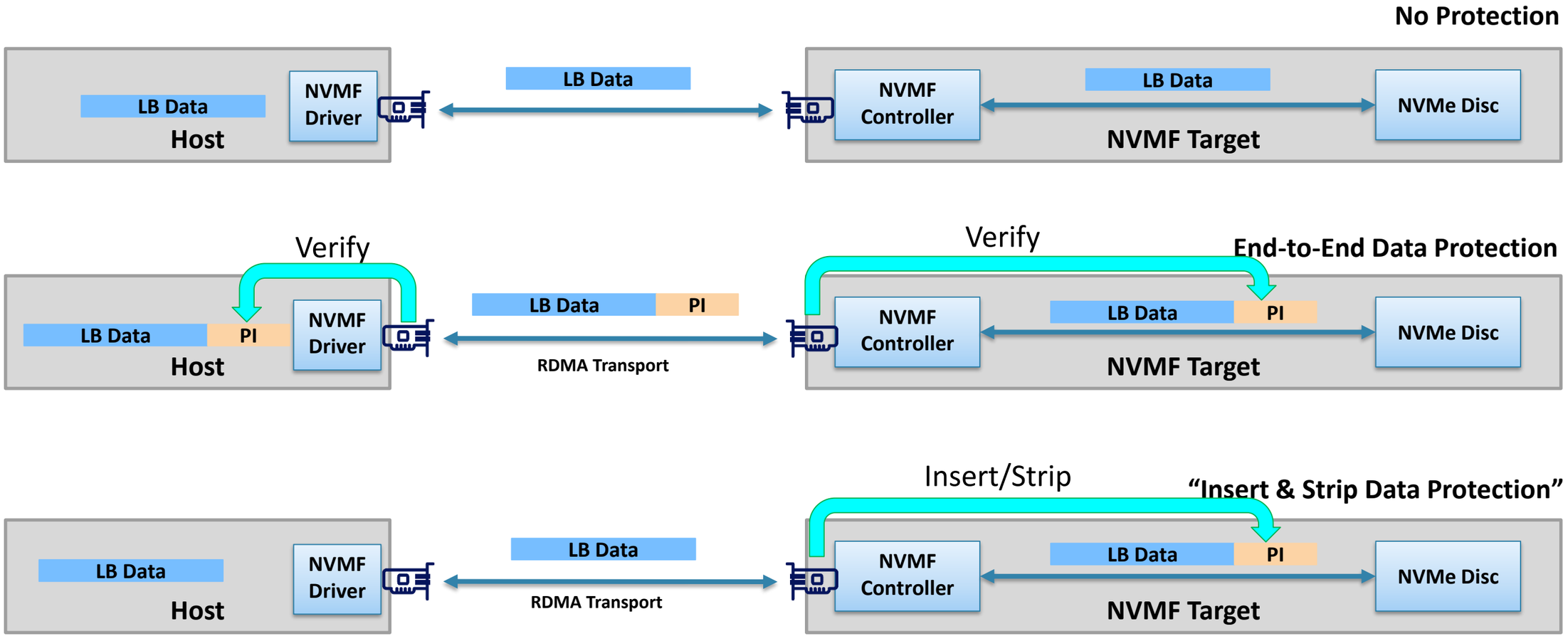
  - Separate: Integrity payload fields lie in a separate buffer from the data.
    - Not supported in Fabrics by definition of the spec (not enough space in the SQE for metadata pointer)



NVME-PCI. Data protection



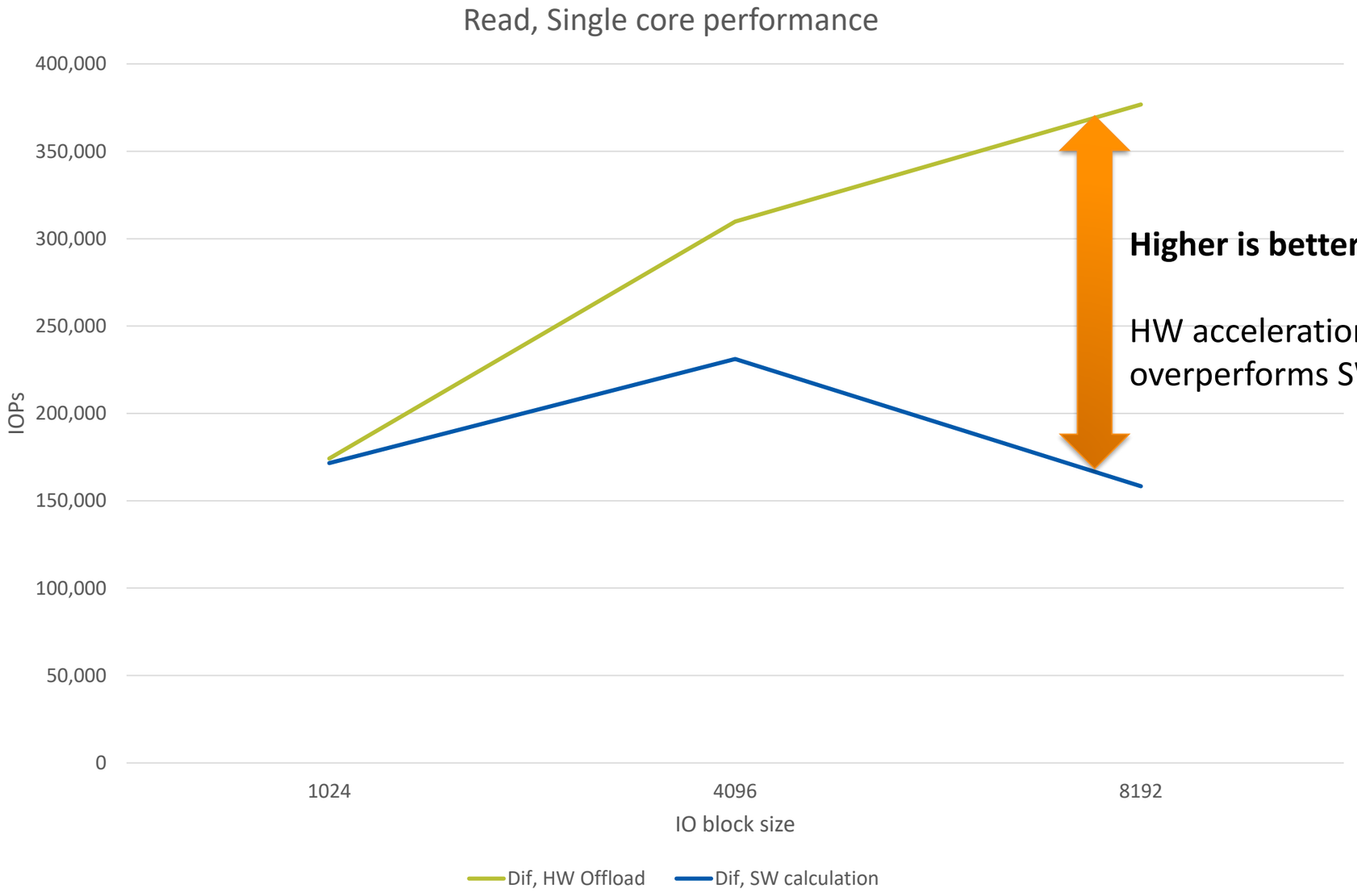
NVME-OF. Data protection



SPDK. DIF “Insert & Strip” mode

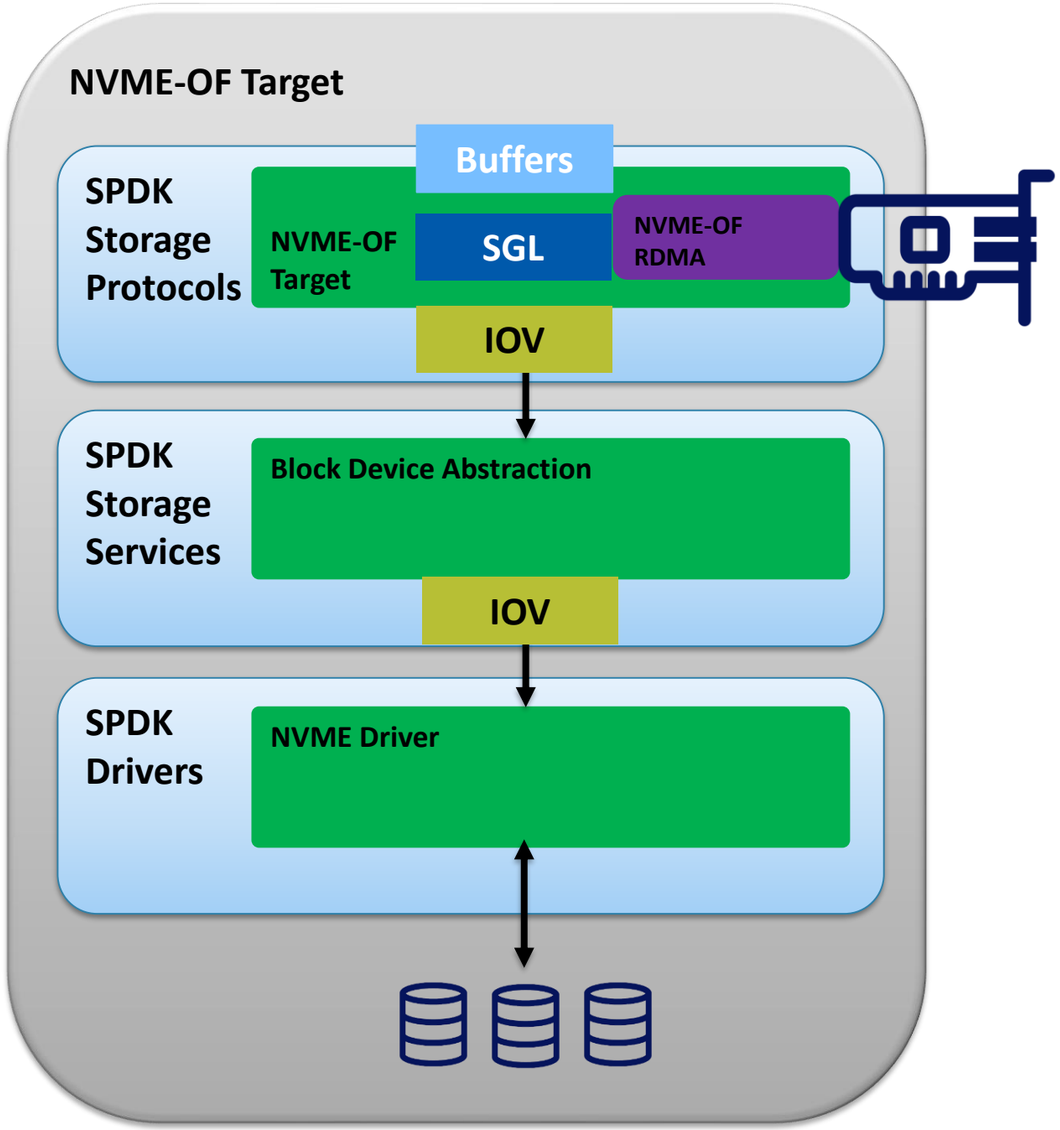
- DIF “Insert & Strip” mode in TCP Transport
  - Shuhei Matsumoto, “Hitachi” : https://review.gerrithub.io/c/spdk/spdk/+/456452 - SW implementation
    - Available in SPDK v19.07
- DIF “Insert & Strip” mode in RDMA Transport
  - Aleksey Marchuk, Evgeny Kochetov, “Mellanox” : https://review.gerrithub.io/c/spdk/spdk/+/465248 - SW implementation
  - HW accelerated mode is under development : https://github.com/EugeneKochetov/spdk/tree/nvmf_rdma_sig_offload

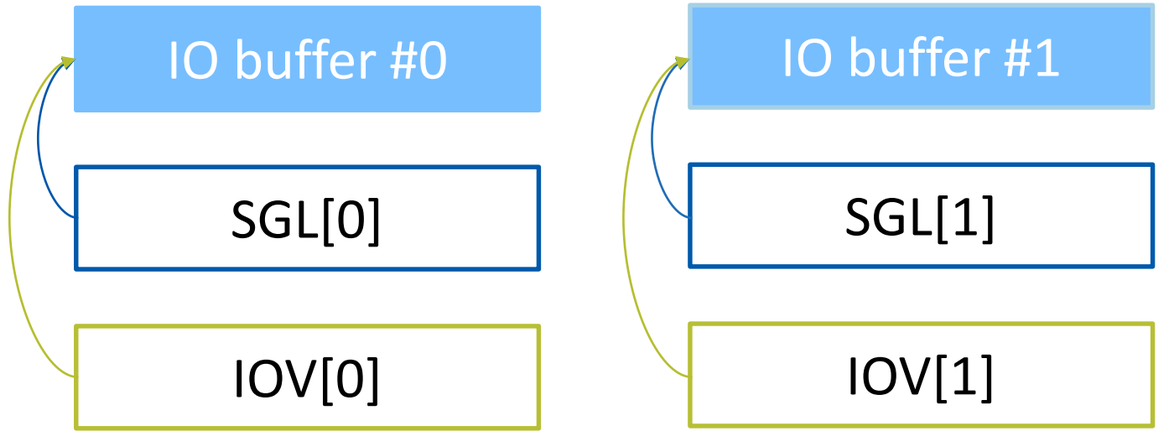
DIF “Insert & Strip” mode. SW vs HW



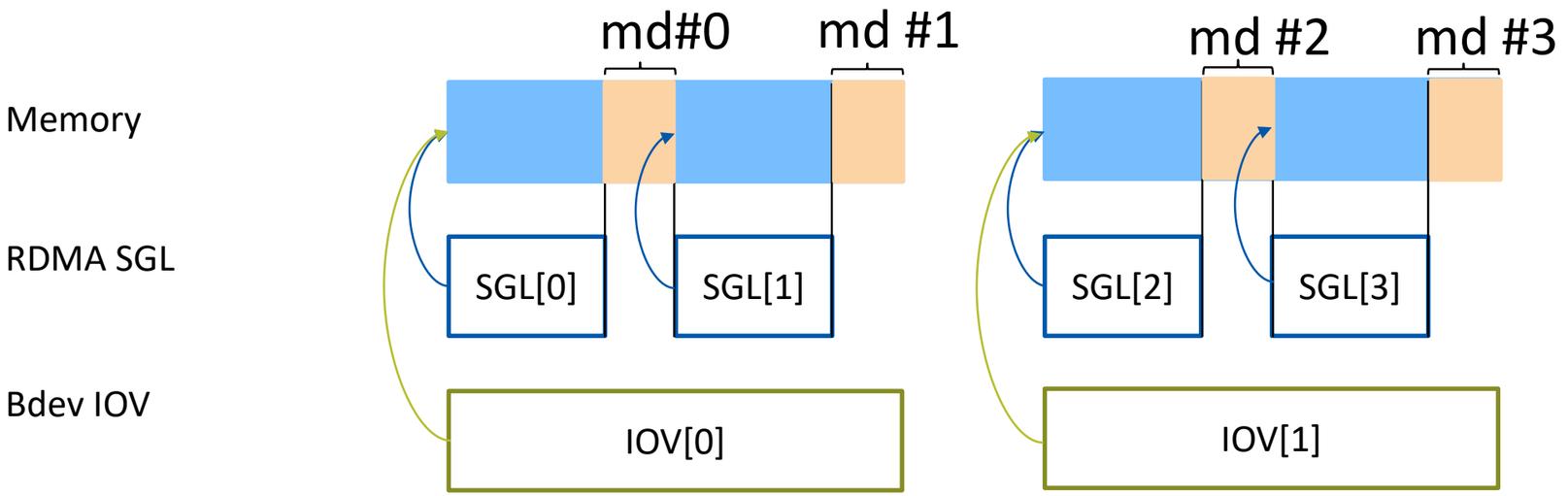
Queue depth: 32
Block size: 512+8
Disk: Samsung PM1725b
Platform: x86

SPDK. Memory management in NVME-OF RDMA

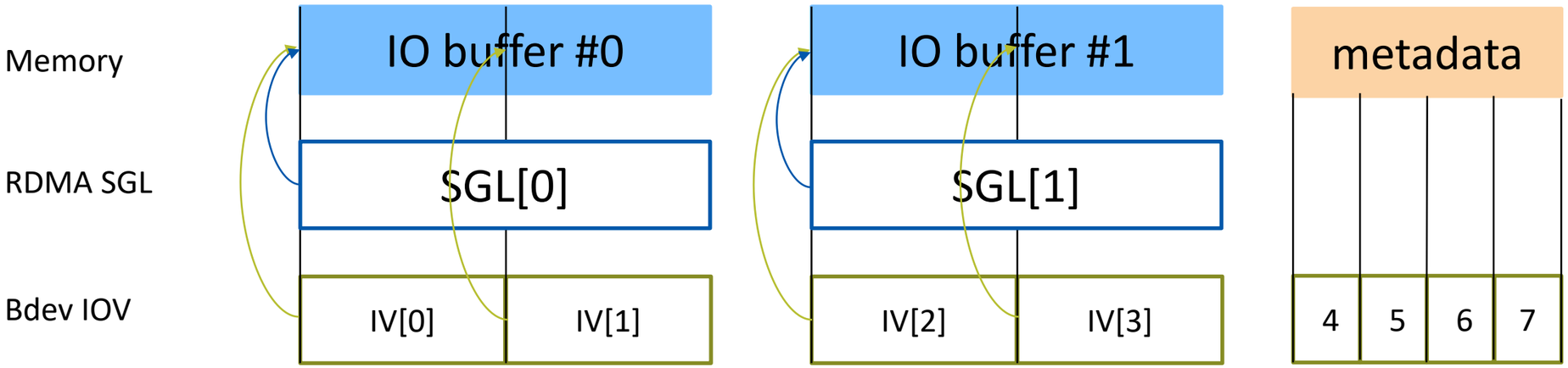




NVME-OF RDMA. Metadata placement



DIF Model

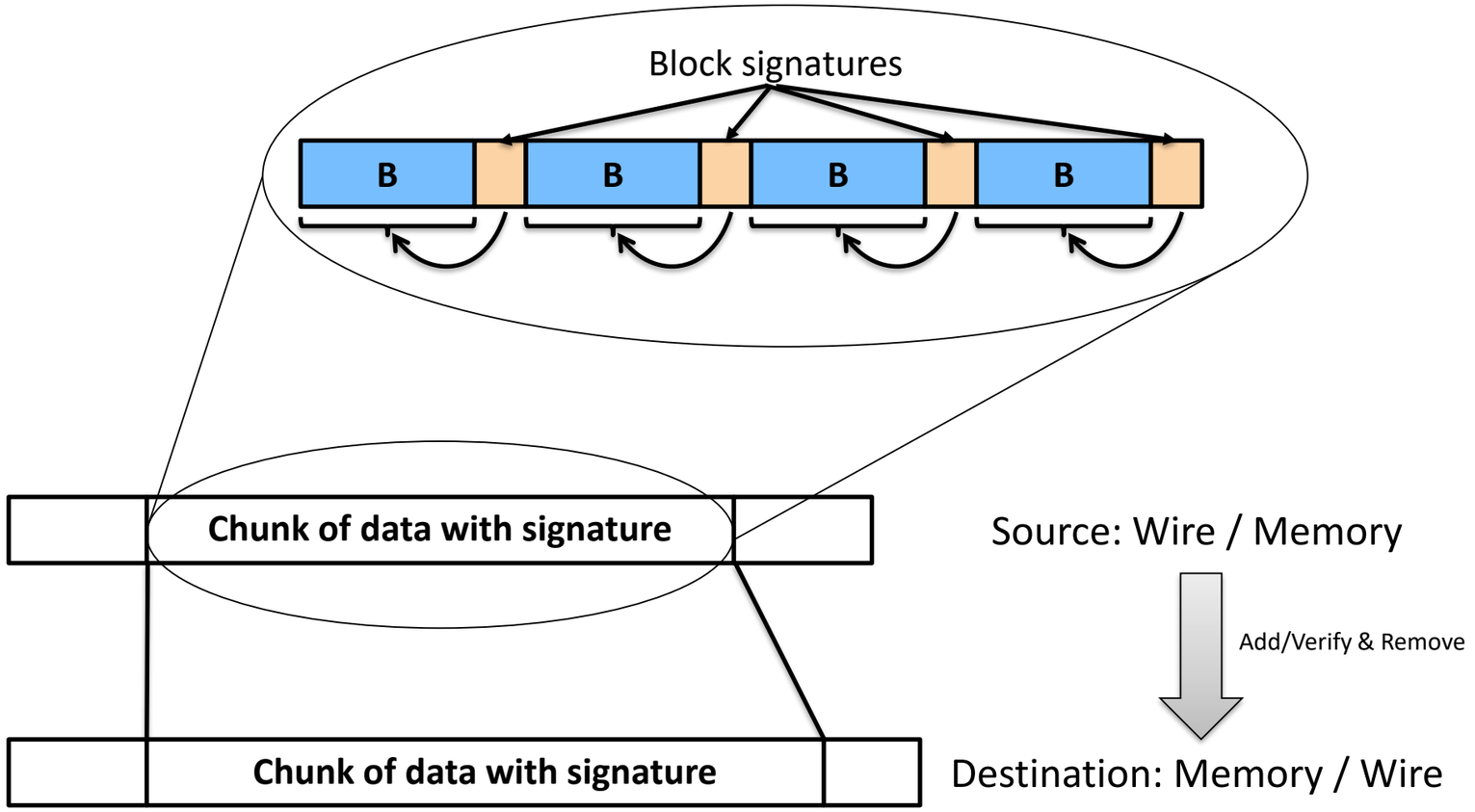


DIX Model

NVME-OF RDMA. Metadata placement

- “DIF” model increases number of SGL elements in RDMA layer
- “DIX” model increases number of IOV elements transferred to bdev layer
- In performance testing “DIF” model overperforms “DIX”
- “DIF” model is chosen as default option
  - Multi-element SGL will be can be replaced by UMR (“User memory region”)
- “DIX” model is used for “in-capsule” data

HW acceleration for “DIF”



“User space” API for “DIF”

- Signature operation is executed at data moving between two Signature Domains
  - Wire Domain
  - Memory Domain
- Signature Operations
  - Add
  - Verify
  - Verify & Remove
- Signature types
  - Repeating block signature. All blocks must have equal size
  - Transaction signature are used for protecting entire transaction
  - Variable block signature covers data of any size
- Using “indirect” memory referencing, both DIF and DIX modes are supported
- Planned to be submitted to “upstream” (rdma-core) in 2019

HW acceleration for data protection. Summary

- HW acceleration for guard tag calculation by NIC demonstrates advantage over SW implementation
- Roadmap:
  - User-spaces API for “DIF” manipulation. Submitting to “upstream”
  - HW acceleration for “Insert & strip” mode in SPDK’s implementation for NVME-OF target
  - HW acceleration for Data Integrity Field generation in SPDK’s initiator
  - Verifying DIF in network layer (RDMA) in “initiator” and “target” sides

Advanced hardware accelerations

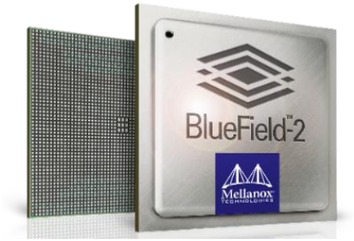
BlueField-2

Superior Storage Performance

- 8 Arm® A72 CPUs @ 2GHz-2.5GHz
- Dual 100Gb/s or Single 200Gb/s ports
- 16 lanes of PCIe Gen4.0
- Up to 5.4M IOPs @ 4KB
- Lowest latency

Storage Accelerations

- NVMe-oF offloads
- NVMe-oF SPDK offload
- RAID, Erasure Coding, CRC32, CRC64 and T10-Diff



Storage Security

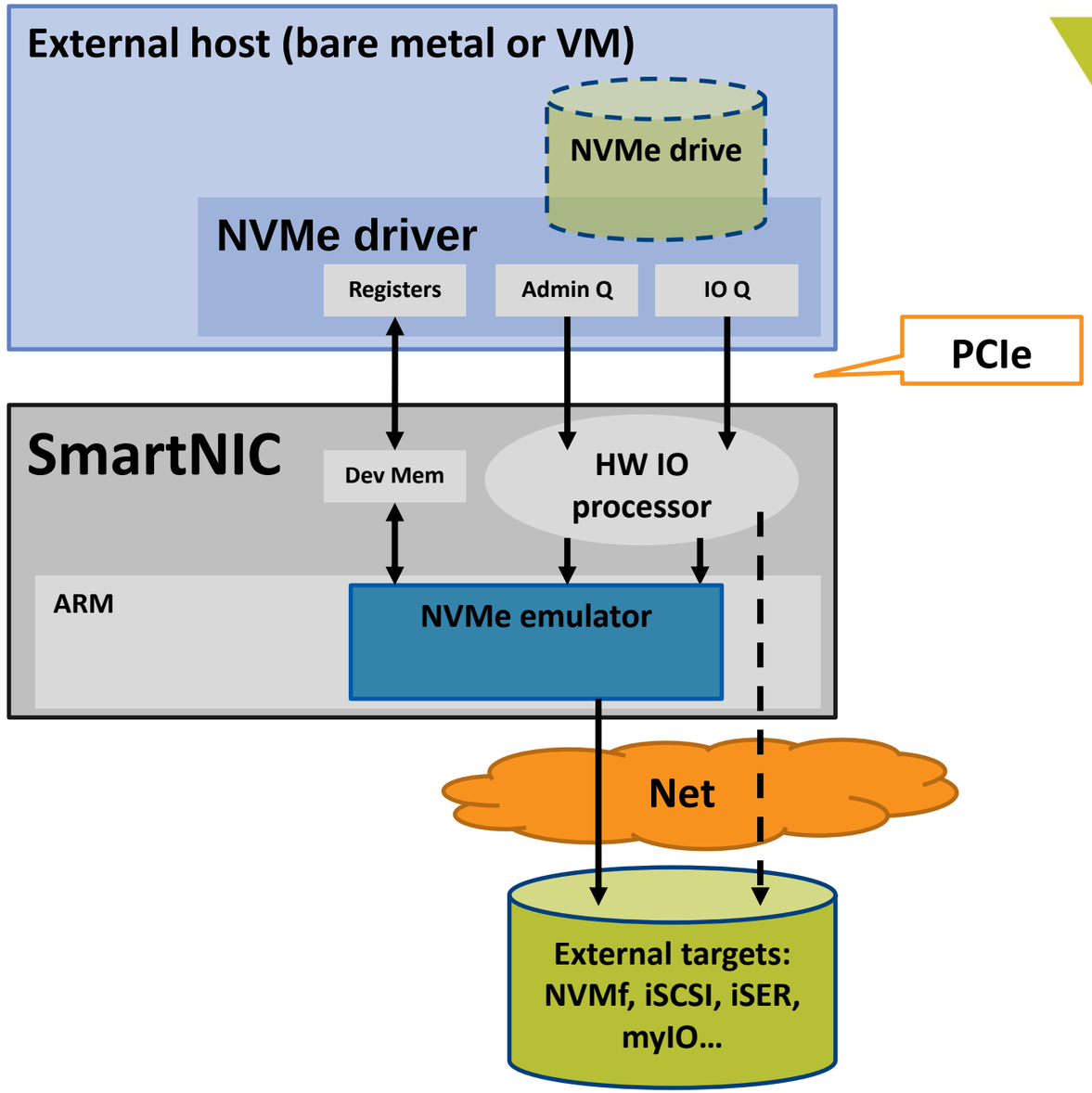
- Data-at-Rest AES-XTS encryption
- Authentication/Authorization services
- Encryption and decryption of data to/from storage
- Protection between users

Unique Features

- Data (De)Compression
- NVMe SNAP™
- Deduplication

NVMe SNAP

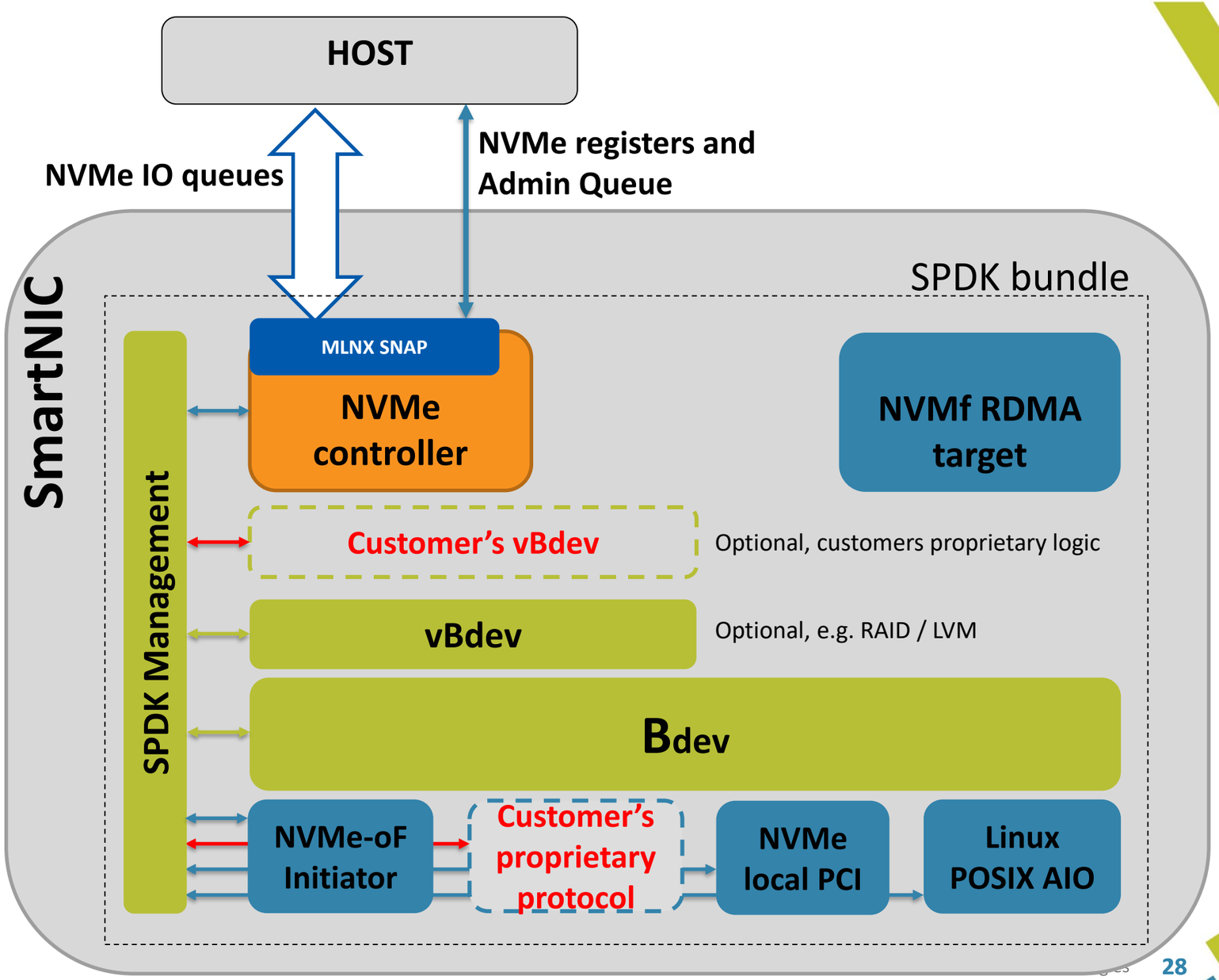
- Emulate locally attached PCIe NVMe drive
- Unmodified NVMe driver on host
- NVMe queues serviced in ARM
  - Then go to network
  - Admin Queue, IO Queues
- Optional: IO path skips ARM
  - Protocol conversion on IO processor
  - Must be simple enough
  - Must be RDMA
  - For example: NVMe-oF
  - Lose IOP-level software manipulation option
  - Admin queue still in ARM



SPDK as NVMe emulators standard framework

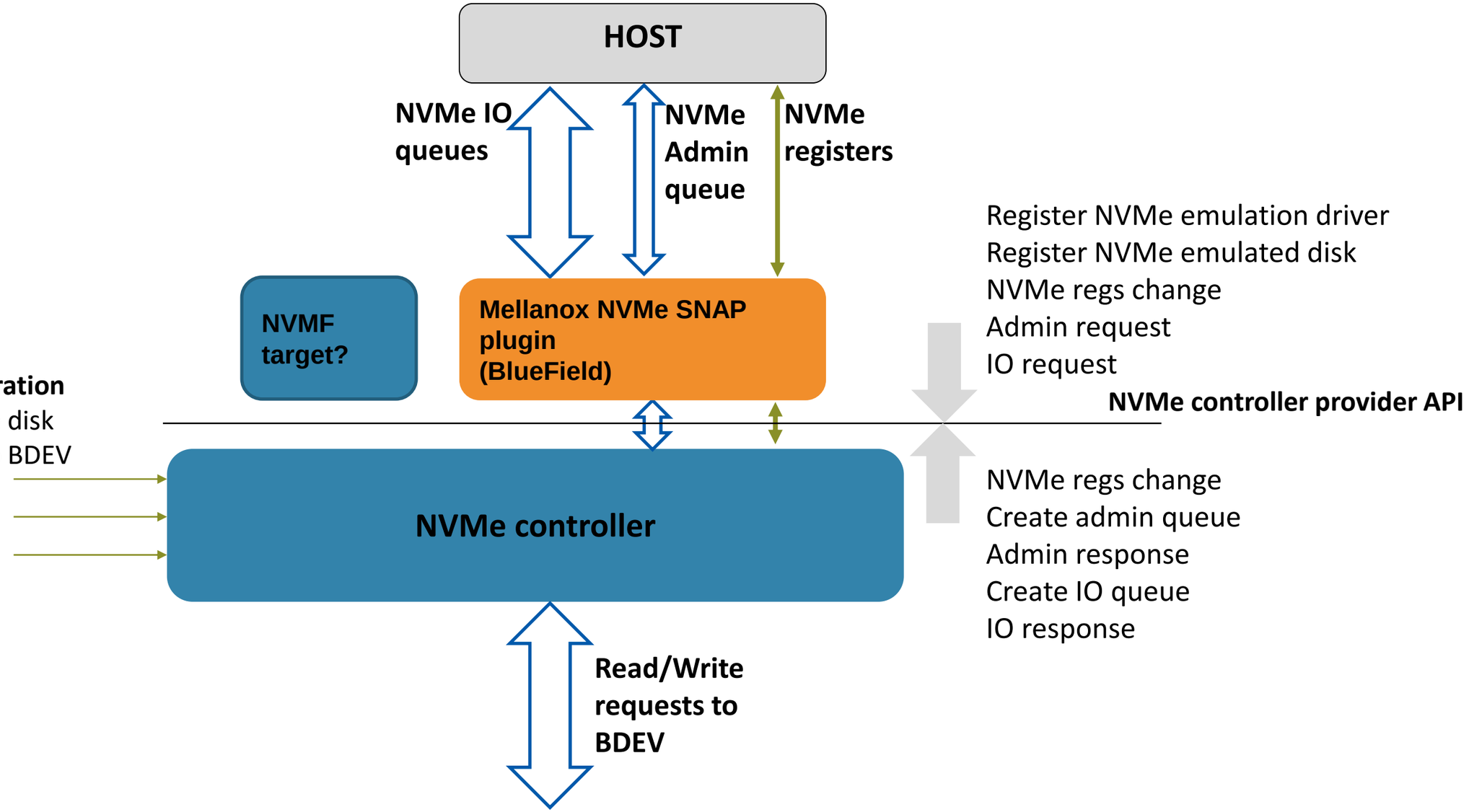
NVMe controller

- New: NVMe controller
  - NVMe device-side registers
  - NVMe device-side admin commands
  - NVMe device-side IO commands
- Vendor specific library
  - Bind to host NVMe device emulation
- Shared code and .h files
  - With NVMe driver
  - With NVMf target
- Configuration is similar to NVMf target
  - Subsystem == emulated NVMe drive
  - Bind BDEVs as Namespaces



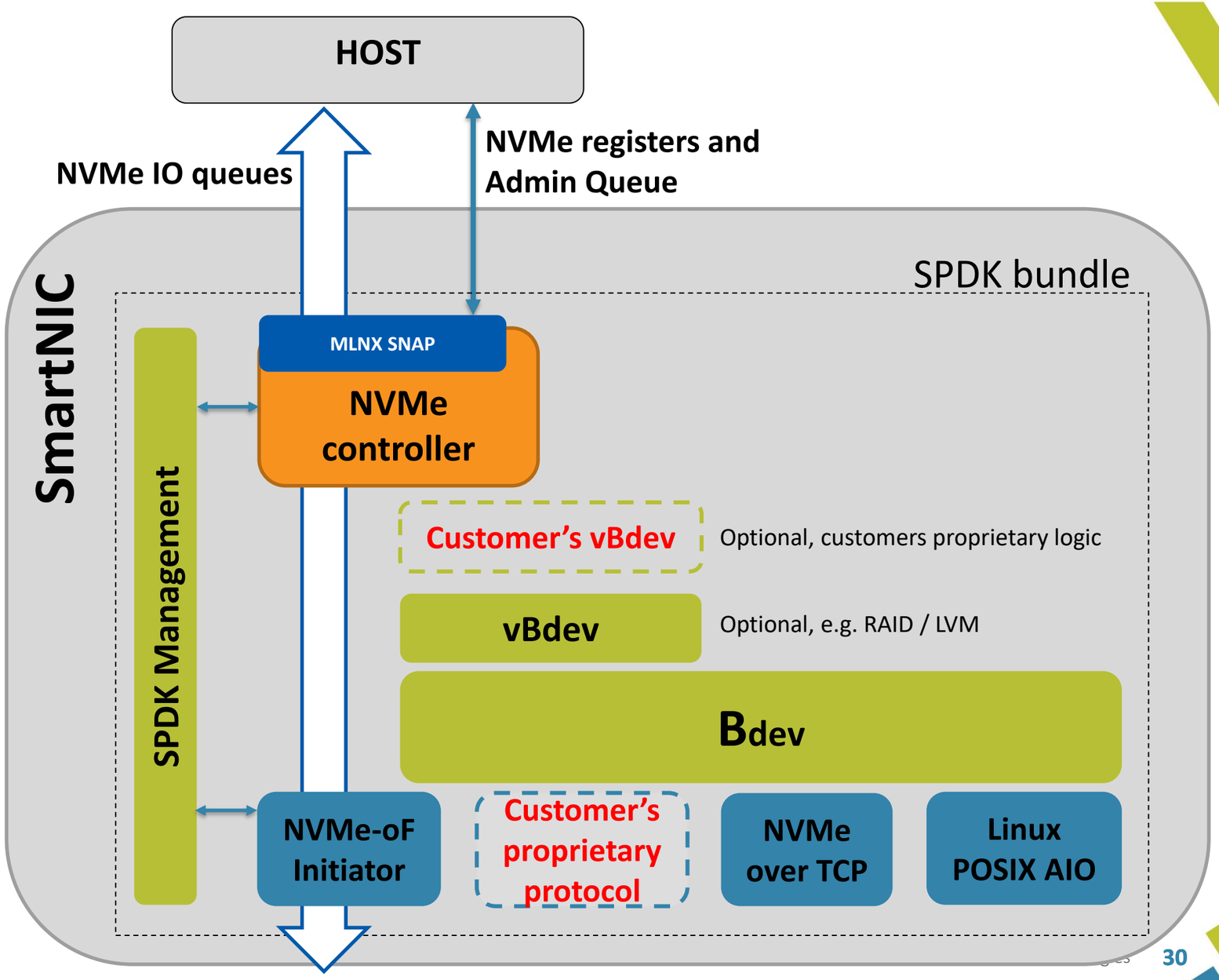
SPDK NVMe Controller

Zoom in



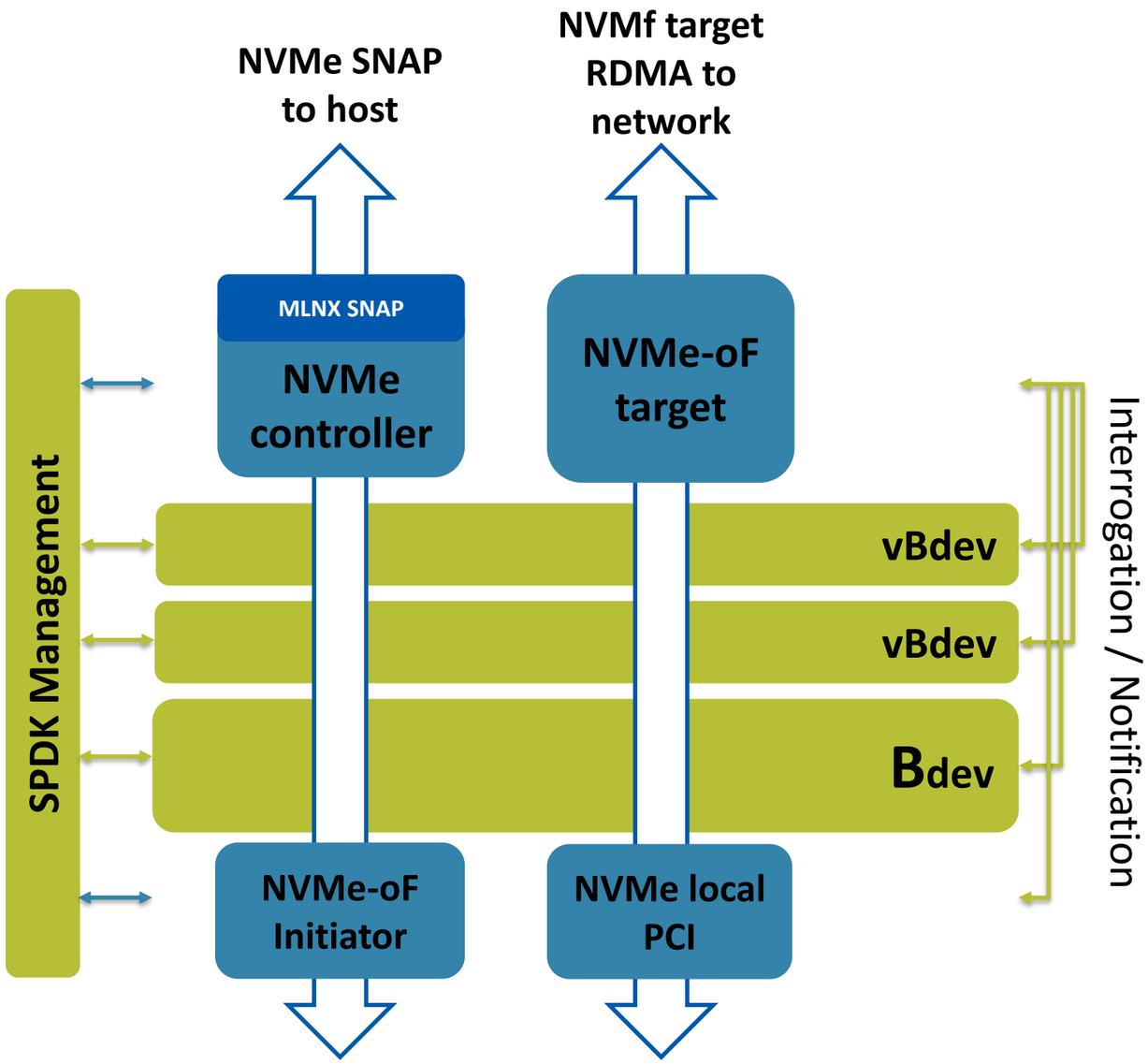
NVMe Controller full-path offload

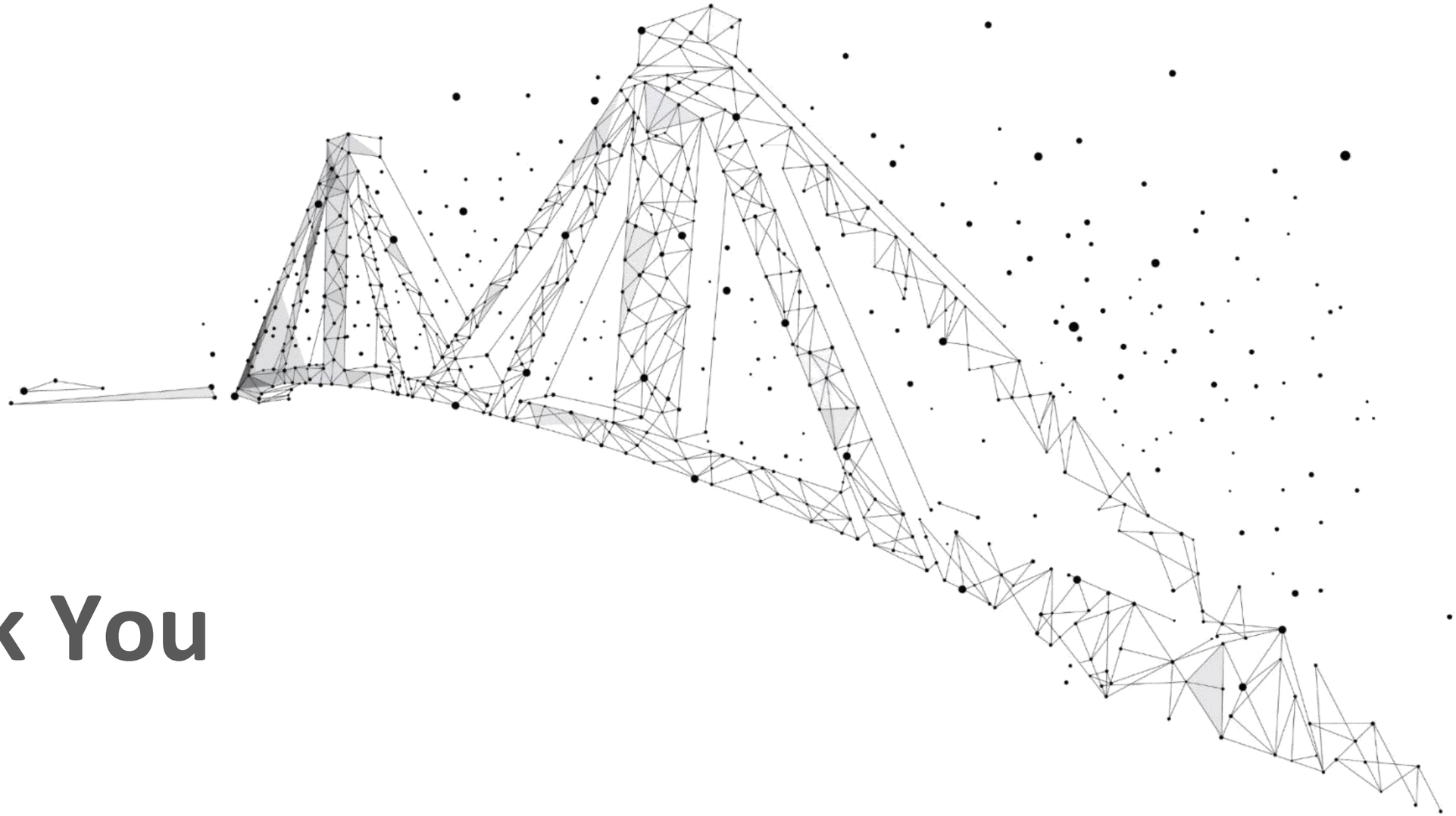
- NVMe SNAP to NVMf initiator offload
- Per emulated device configuration
  - Don’t offload
  - Always offload
    - Fail configuration if not possible
  - Best effort offload
    - Offload if possible, software path if not
- Best performance!
  - For simple use cases



SPDK in-network offloads

- vs. local mem-to-mem offloads
- Upper application configured to use a **bdev**
  - NVMe controller for SNAP
  - NVMe-oF target
- Interrogate vbdevs/bdevs chain
  - Identify the kind of bdev (NVMf, iSCSI, Crypto…)
  - Get configuration / create resources
  - If vbdev, get next (v)bdev(s), repeat
- Can the full flow and configuration be offloaded?
  - If yes – allow offload, configure device
  - If no – continue in software
- Notification for runtime changes in configs
  - Thin provisioning new chunk mapped
  - Volume resized

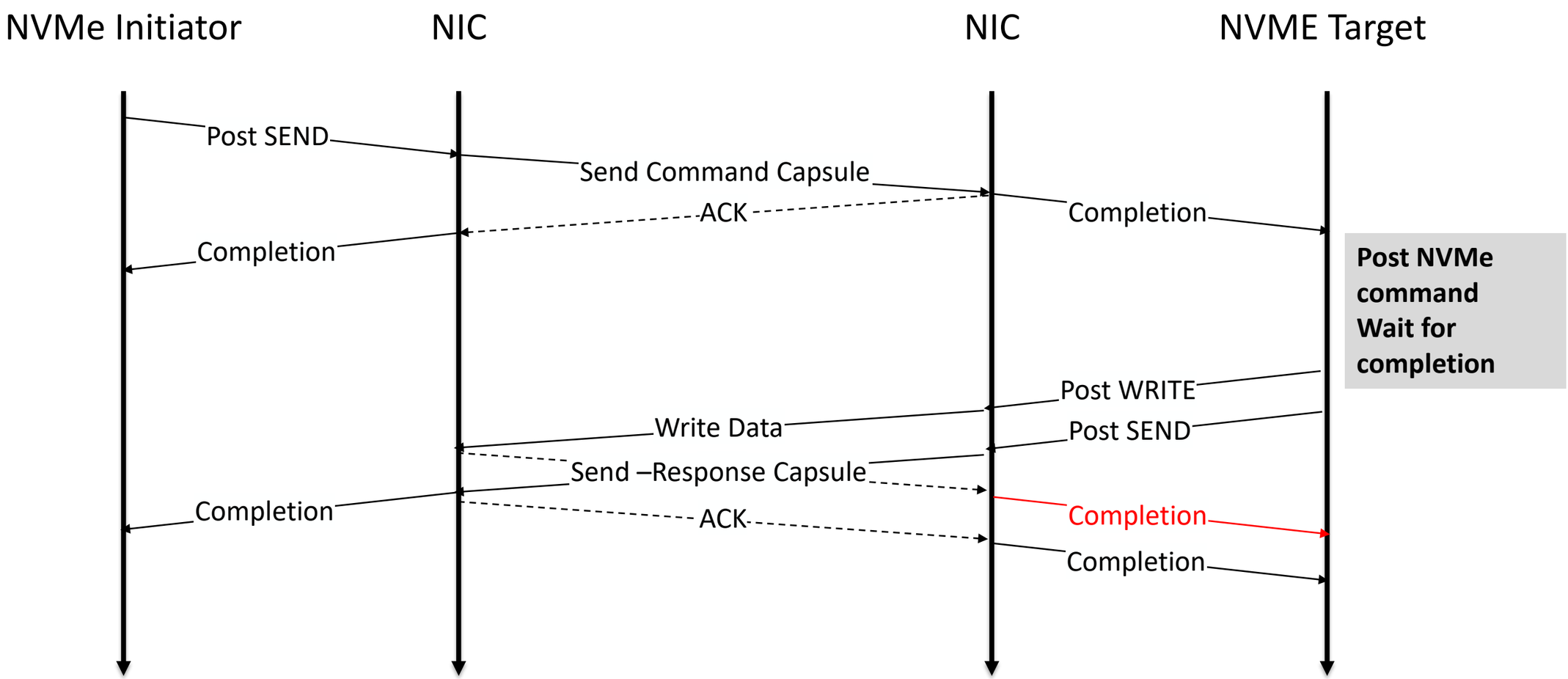




Thank You

Backup

NVME-OF RDMA. IO Read. Selective signaling



NVME-OF RDMA. Request batching