

PERSISTENT MEMORY PROGRAMMING MADE EASY WITH PMEMKV

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AGENDA

Why pmemkv?

- Persistent Memory programming is difficult
- key-value store

pmemkv design

- goals for pmemkv
- architecture
- configuration
- life-cycle (persistent libpmemobj-based engines)

Engines

- overview
- multiple engines within the same memory pool
 Language bindings
- pmemkv is simple!
 - API
 - C++ example
 - NodeJS example
- Latencies and performance Q&A

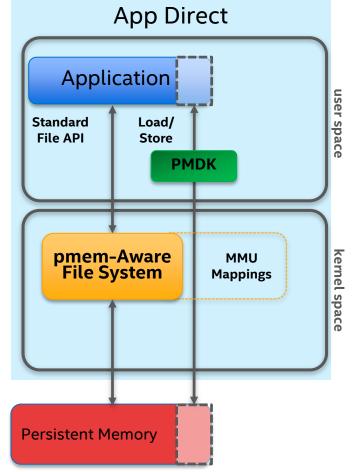


WHY PMEMKV?





Why pmemkv? Persistent Memory programming is difficult

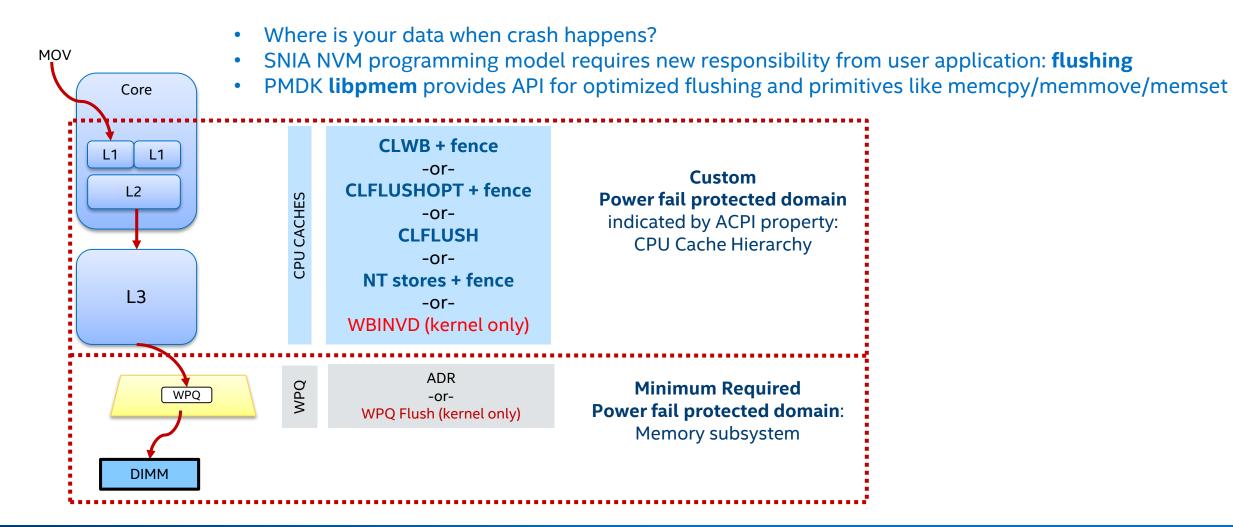


- Different modes for using Persistent Memory:
 - Memory Mode
 - Storage over App Direct
 - App Direct
- A pmem-aware file system exposes persistent memory to applications as files
- In-place persistence (no paging, context switching, interrupts, nor kernel code executes)
- Byte addressable like memory (Load/store access, no page caching)
- Cache Coherent

It looks easy! But...



Why pmemkv? Persistent Memory programming is difficult

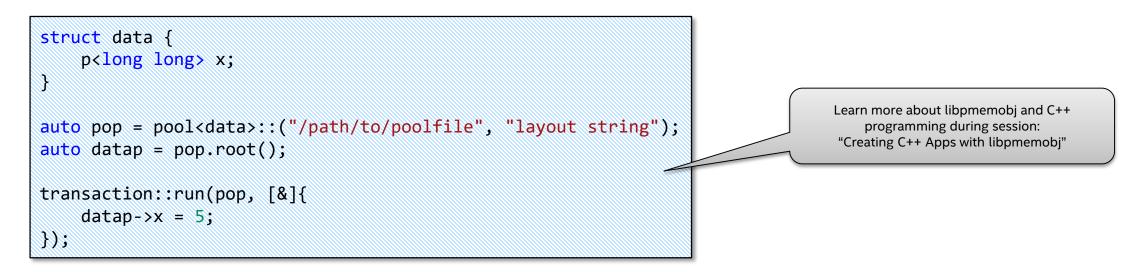


Why pmemkv? Persistent Memory programming is difficult



Result?

- 1. "0 0 0 0 0 0 0 0 0 0 0 ..."
- 2. "Hello, W\0\0\0\0\0..."
- 3. "\0\0\0\0\0\0\0orld!\0"
- 4. "Hello, \0\0\0\0\0\0\0
- 5. "Hello, World!\0"
- pmem_persist is faster than msync(), but it is still not transactional
- SNIA NVM programming model requires new responsibility from user application: consistency
- PMDK libpmemobj provides transactional API, Persistent Memory allocator etc.



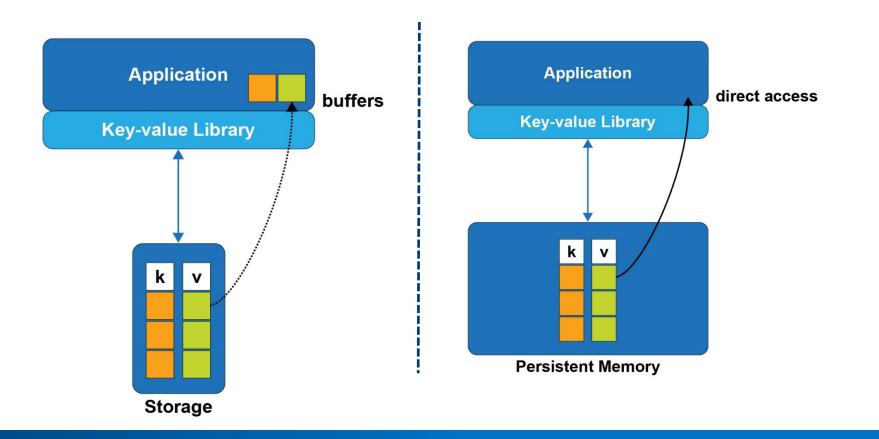
Why pmemkv?

- Q: How to make Persistent Memory programming easier?
- A: Local Key-Value data store
- API flexibility increases complexity
 - API flexibility not always desired
- Usually the bigger barrier to adoption, the better performance gains
 - Don't have to be true for some specific workloads
- Large addressable market of cloud developers for an easy KV store
 - Data stored in cloud will be the majority of all stored data in nearby future
- Key-Value data store provides straightforward API which can easily utilize Persistent Memory advantages
 - Nothing new to learn in order to start using Persistent Memory in efficient way
- Simple API makes creation of different language bindings relatively easy
 - Important in cloud native computing, where many high-level languages are being used



Why pmemkv?

- Key-value store can take advantage from persistence and big capacity of Persistent Memory
- Key-value store can utilize Persistent Memory byte addressability
 - huge performance gain for relatively small key and values



PMEMKV DESIGN



pmemkv design goals for pmemkv

Technical:

- Local key/value store (no networking)
- Idiomatic language bindings
- Simple, familiar, bulletproof API
- Easily extended with new engines
- Optimized for persistent memory (limit copying to/from DRAM)
- Flexible configuration, not limited to a single storage algorithm
- Generic tests

Community:

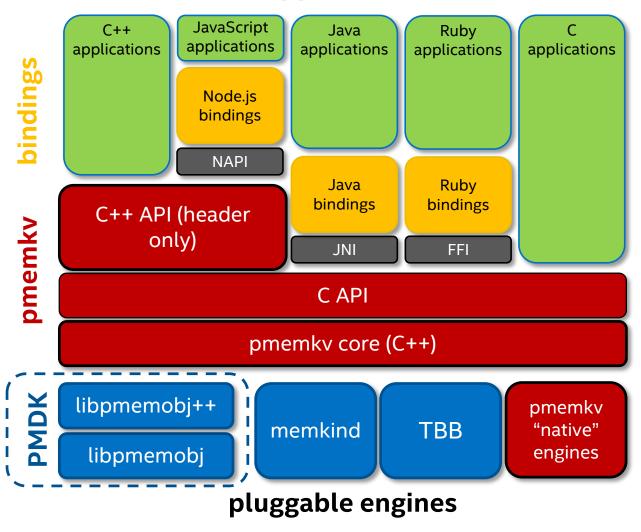
- Open source, developed in the open and friendly licensing
 - https://github.com/pmem/pmemkv
- Outside contributions are welcome
- Intel provides stewardship, validation on real hardware, and code reviews
- Standard/comparable benchmarks



pmemkv design architecture

- pmemkv core is a frontend for engines
 - Core implementation written in C++, not related to Persistent Memory
- Pluggable engines
 - Some engines are implemented in pmemkv, some engines are imported from external projects
 - Persistent engines are implemented with libpmemobj (PMDK)
- Native API for pmemkv is written C/C++
- pmemkv design allows for easy integration with high-level language bindings

applications



pmemkv design configuration

- Flexible configuration API
 - Works with different kinds of engines
- Every engine has documented supported config parameters individually
- Unordered map
 - Takes name configuration value as a k-v pair
- Supported configuration types:
 - int64/uint64/double
 - string
 - Arbitrary data (pointer and size)
- Resides on stack
 - Takes optional destructor as an additional parameter if custom configuration parameter allocates memory

Typical config structure example for libpmemobj-based engines

config cfg;

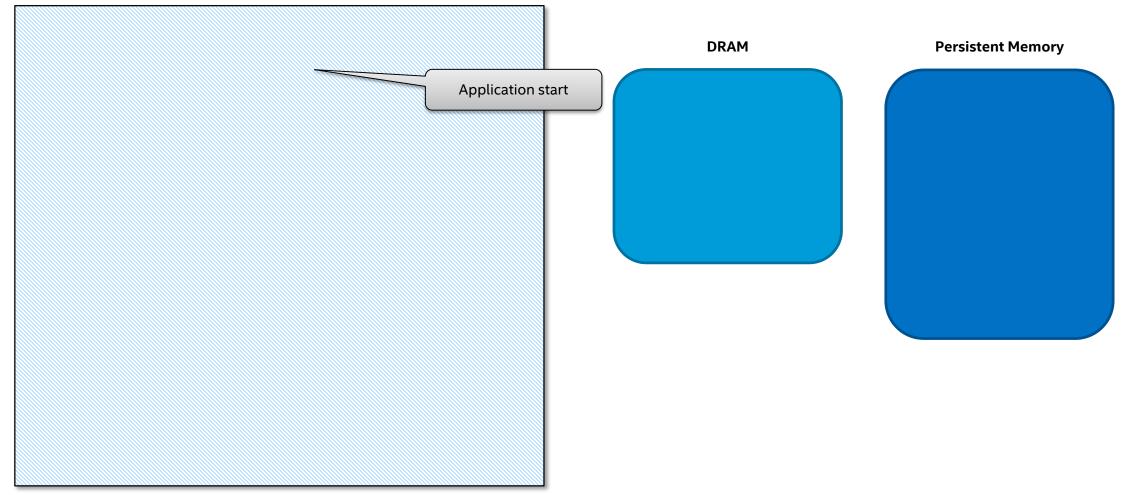
```
status s = cfg.put_string("path", path);
assert(s == status::OK);
```

s = cfg.put_uint64("size", SIZE);
assert(s == status::OK);

```
s = cfg.put_uint64("force_create", 1);
assert(s == status::OK);
```

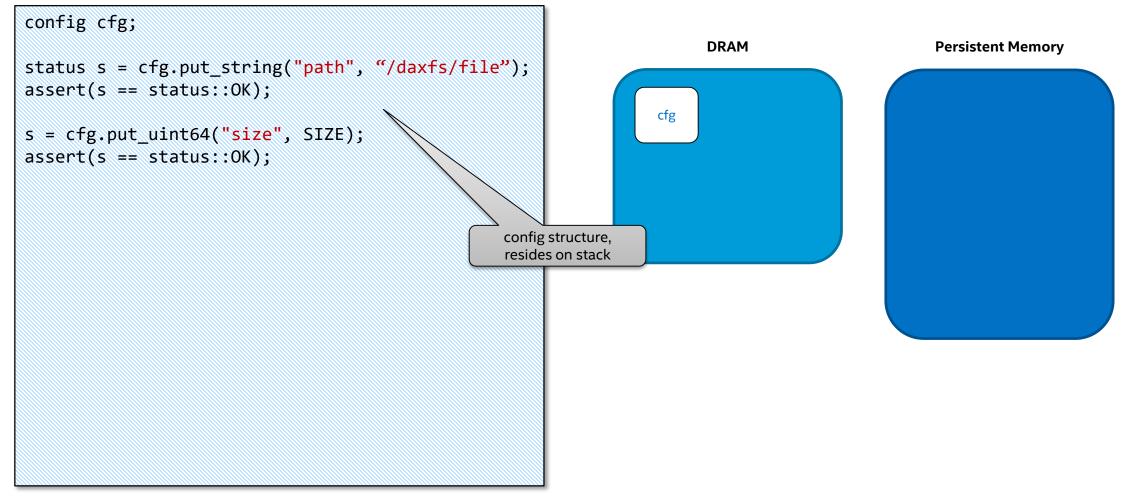


pmemkv design life-cycle (persistent engines based on libpmemobj)



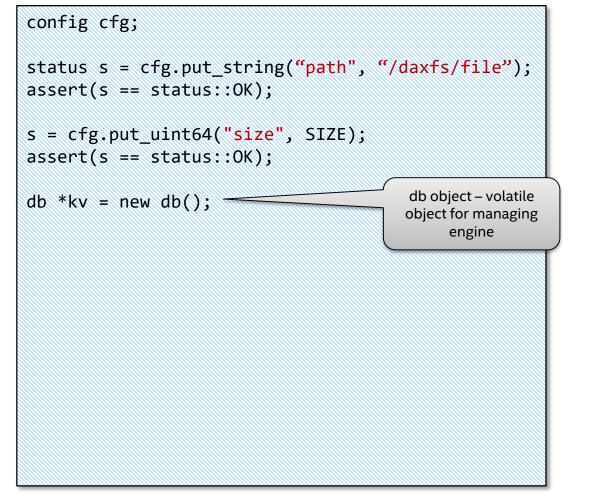


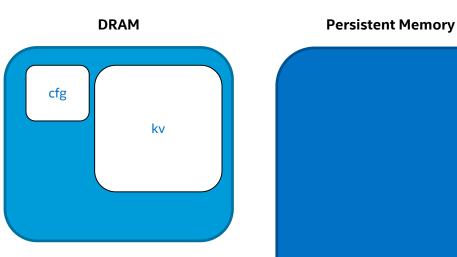
life-cycle (persistent engines based on libpmemobj)



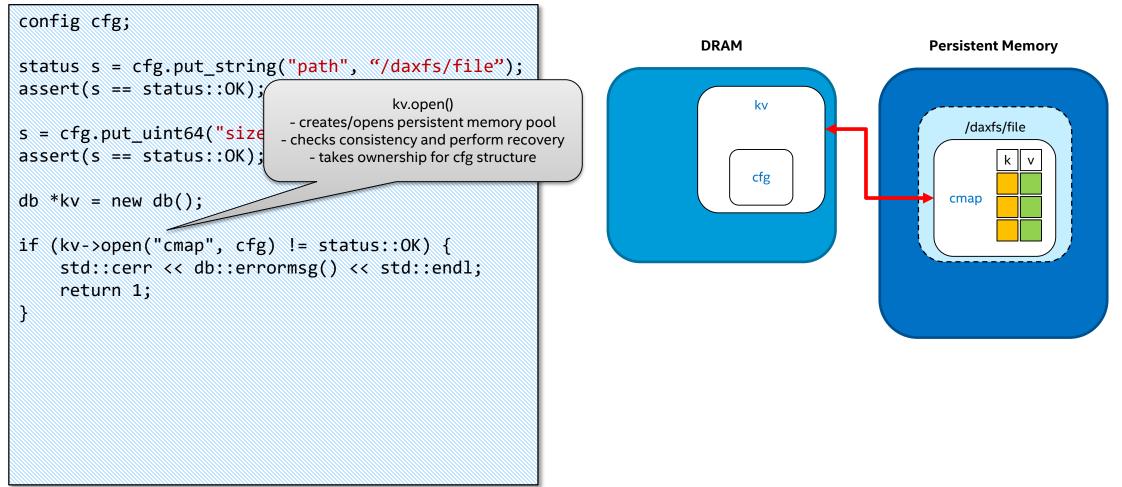


life-cycle (persistent engines based on libpmemobj)





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life-cycle (persistent engines based on libpmemobj)

```
config cfg;
```

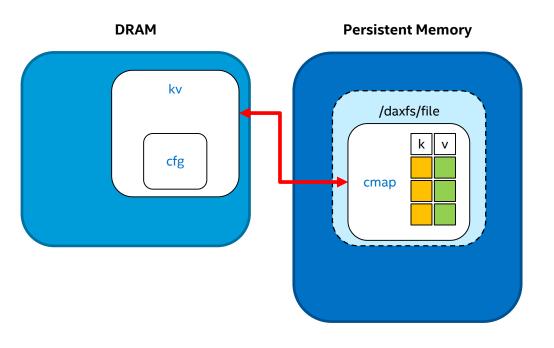
```
status s = cfg.put_string("path", "/daxfs/file");
assert(s == status::OK);
```

```
s = cfg.put_uint64("size", SIZE);
assert(s == status::OK);
```

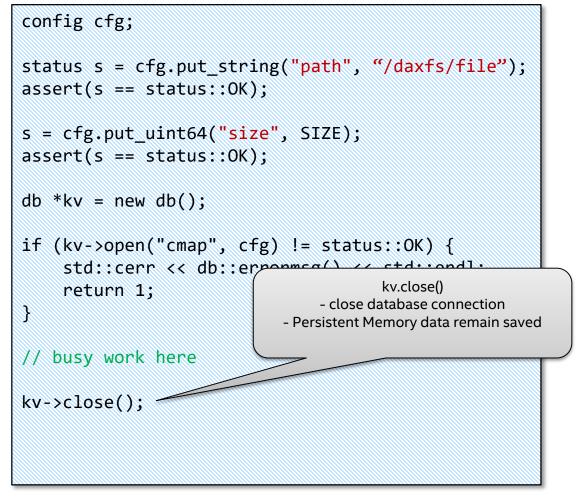
```
db *kv = new db();
```

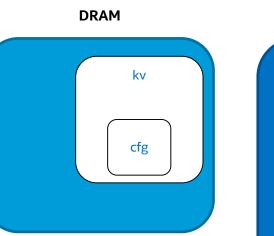
```
if (kv->open("cmap", cfg) != status::OK) {
   std::cerr << db::errormsg() << std::endl;
   return 1;</pre>
```

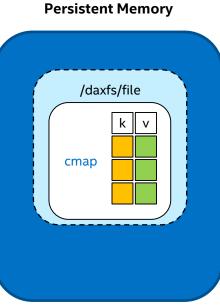
// busy work here



life-cycle (persistent engines based on libpmemobj)



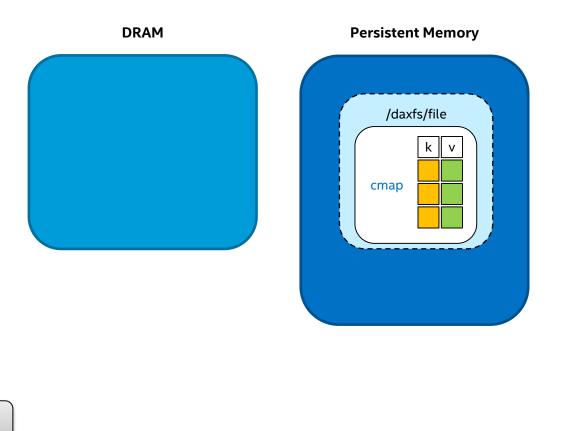






life-cycle (persistent engines based on libpmemobj)

```
config cfg;
status s = cfg.put_string("path", "/daxfs/file");
assert(s == status::OK);
s = cfg.put uint64("size", SIZE);
assert(s == status::OK);
db *kv = new db();
if (kv->open("cmap", cfg) != status::OK) {
    std::cerr << db::errormsg() << std::endl;</pre>
    return 1;
// busy work here
kv->close();
                                Safety deletion of volatile data
delete kv;
```







Engines overview

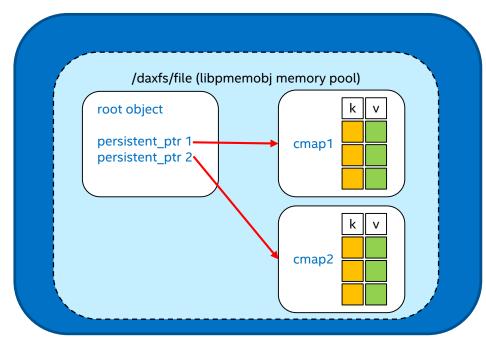
- Engine contributions are welcome!
- Types:
 - ordered/unordered
 - persistent/volatile
 - concurrent/single threaded
- Engines are optimized for different workloads & capabilities
- All engines work with all language bindings
- Generic tests for engines incl:
 - memcheck
 - helgrind/drd
 - pmemcheck
 - pmemreorder

Engine Name	Description	Experimental?	Persistent?	Concurrent?	Sorted?
<u>blackhole</u>	Accepts everything, returns nothing	No	-	-	-
<u>cmap</u>	Concurrent hash map	No	Yes	Yes	No
<u>vsmap</u>	Volatile sorted hash map	No	No	No	Yes
<u>vcmap</u>	Volatile concurrent hash map	No	No	Yes	No
tree3	Persistent B+ tree	Yes	Yes	No	No
<u>stree</u>	Sorted persistent B+ tree	Yes	Yes	No	Yes
<u>caching</u>	Caching for remote Memcached or Redis server	Yes	Yes	No	-
csmap	Sorted concurrent map (under development)	Yes	Yes	Yes	Yes

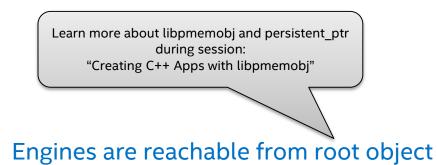


Engines multiple engines within the same memory pool

- pmemkv API (config API) does not limit user for correlating single engine with single memory pool (libpmemobj)
- It is possible to pass persistent_ptr argument to config structure and attach engine to the pointer (ongoning work on generic API for libpmemobj-based engines)



Persistent Memory





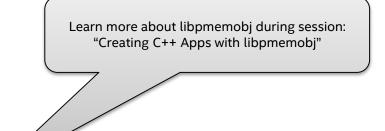
Engines

multiple engines within the same memory pool

```
struct Root {
    pmem::obj::persistent_ptr<PMEMoid> ptr1;
    pmem::obj::persistent_ptr<PMEMoid> ptr2;
};
// libpmemobj setup here
config cfg_1;
config cfg_2;
status ret = cfg_1.put_object("oid", &(pop.root()->ptr1),
nullptr);
assert(ret == status::OK);
ret = cfg_2.put_object("oid", &(pop.root()->ptr2), nullptr);
assert(ret == status::OK);
```

```
db *kv_1 = new db();
status s = kv_1->open("cmap", std::move(cfg_1));
assert(s == status::OK);
```

```
db *kv_2 = new db();
s = kv_2->open("cmap", std::move(cfg_2));
assert(s == status::OK);
```



Prototyped API for using pmemkv with libpmemobj++ simultaneously (implementation work ongoing)



LANGUAGE BINDINGS



Language bindings

Simple API = easy to implement high-level language bindings with small performance overhead

- Currently 4 available language bindings for pmemkv:
 - Java <u>https://github.com/pmem/pmemkv-java</u>
 - NodeJS <u>https://github.com/pmem/pmemkv-nodejs</u>
 - Ruby <u>https://github.com/pmem/pmemkv-ruby</u>
 - Python https://github.com/pmem/pmemkv-python
- Their APIs are not functionally equal to native C/C++ counterpart
 - Configuration possible only by JSON string passed to open() function
 - Multiple engines within single memory pool not possible
 - Above API gaps are under development



PMEMKV IS SIMPLE!



pmemkv is simple!

- Well understood key-value API
 - Nothing new to learn
 - Inspired by rocksDB and levelDB
- Life-cycle API
 - open()/close()
- Operations API
 - put(key, value)
 - get(key, value/v_callback)
 - remove(key)
 - exists(key)

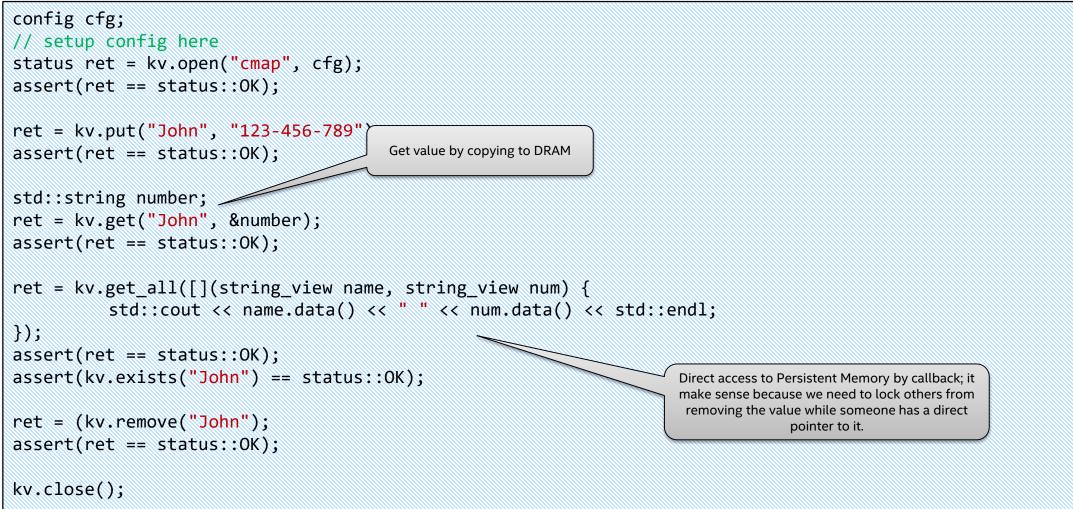
- other
 - errormsg()
- Iteration API
 - count_all()
 - get_all(kv_callback)
- +range versions of above for ordered engines
 - below/above/between

pmemkv is not limited to the API above – in future, specific engines might provide extensions and methods like batch_update()



pmemkv is simple!

C++ example



pmemkv is simple! NodeJS example

```
const db = new Database('cmap', '{"path":"/daxfs/kvfile","size":1073741824}');
db.put('John', '123-456-789');
assert(db.get('John') === '123-456-789');
db.get_all((k, v) => console.log(`name: ${k}, number: ${v}`));
db.remove('John');
assert(!db.exists('John'));
db.stop();
```

• Similar simplicity with other high-level language bindings

LATENCIES AND PERFORMANCE



Latencies and performance

- Language bindings
 - number of round trips between high-level language & native code
 - Create high-level object (string, byte[], reference type, callback/closure)
 - Translate bytes to UTF-8
 - String interning, reference counting or GC
- pmemkv core (native code)
 - Searching indexes in DRAM
 - Updating indexes in DRAM
 - Managing transactions
 - Allocating persistent memory
- Persistent Memory
 - HW read and write latency
- Performance varies based on traffic pattern
 - Contiguous 4 cacheline (256B) granularity vs. single random cacheline (64B) granularity
 - Read vs. writes

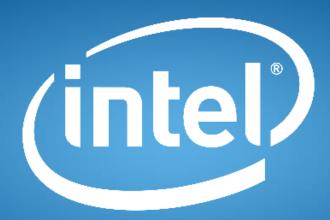
Latencies and performance cmap performance

- pmemkv_tools is a separate github repository with benchmark tool inspired by db_bnch
 - https://github.com/pmem/pmemkv-tools

- Test results for cmap (persistent concurrent hashmap)
 - Throughput scales with a number of threads
 - P99 latency flat







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